**Connectivity concepts of an Arithmetic Graph**

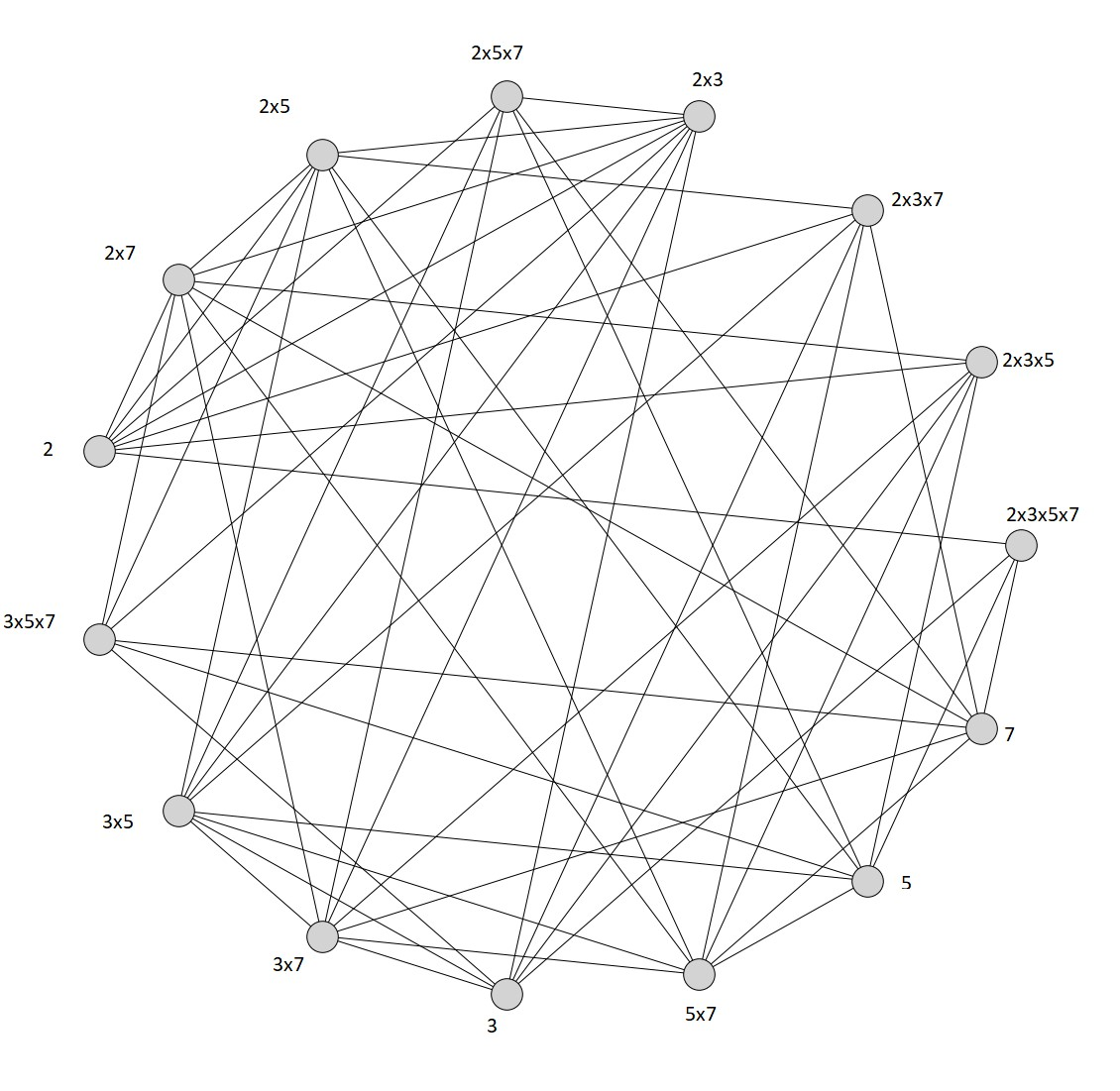
**1.Introduction**

Number theory is one of the oldest branches of mathematics which inherited rich contributions from almost all greatest mathematicians ancient and modern. The theory of congruences in Graph theory by Melvyn Bernard Nathanson in [3], paved the way for the emergence of a new class of graphs, namely Arithmetic Graphs. Inspired by the interplay between number theory and graph theory several researches in recent times are carrying out extensive studies on various arithmetic graphs in which adjacency between vertices is defined through various arithmetic functions. Vasumathi and Vangipuram defined the arithmetic graph in such a way that the adjacency between the vertices of same parity is considered in [6]. Suriyanarayana Rao and Sreenivasan.V in [5] defined the arithmetic graph by excluding the condition of adjacency between the vertices of same parity. By the fundamental theorem of arithmetic, every positive integer greater than one can be uniquely represented as a product of primes (i.e) where *ai* ≥ 1 . Here we discuss about the positive integer other than where *p* is prime and *ai*≥ 1 .

**2**.**Arithmetic graph**

**Definition2.1.** The **arithmetic graph Vn** is defined as a graph with its vertex set is the set consists of the divisors of n (excluding 1) where n is a positive integer and where *pi* ′ s are distinct primes and *ai* ′ s ≥ 1 and two distinct vertices *a, b* which are not of the same parity are adjacent in this graph if (*a, b*) = *pi* , for some *i* , 1 ≤ i ≤ r. The vertices *u* and *v* are said to be of same parity if both *u* and *v* are the powers of the same prime, for instance *u* = *p* 2 *, v* = *p* 3 . In this graph, vertex 1 becomes an isolated vertex. Hence, we consider arithmetic *V*n graph without vertex 1. Therefore, each vertex of Vn is connected to some vertex in *V*n . Clearly, *V*n is a connected simple graph.

**Example 2.2.** Consider an arithmetic graph G = V210 where 210 = 2 × 3 × 5 × 7. The vertex set V (G) = {2, 3, 5, 7, 2×3, 2×5, 2×7, 3×5, 3×7, 5×7, 2×3×5, 2×3×7, 2× 5×7, 3×5×7, 2×3×5×7}. Since (2, 2×3) = (2, 2×5) = (2, 2×7) = (2, 2×3×5) = (2, 2×3×7) = (2, 2×5×7) = (2, 2×3×5×7) = 2 there exist edges between the vertex 2 and the vertices such as 2×3, 2×5, 2×7, 2×3×5, 2×3×7, 2×5×7, 2×3×5×7. Also, since the (2 × 3, 2 × 3 × 7) = 2 × 3= 6≠ pi , 2 × 3 and 2 × 3 × 7 are non-adjacent vertices. Similarly for the other vertices. This can be shown in Figure 2.1



**Figure 2.1**: **Arithmetic Graph** *G=V*210

**Theorem 2.3.[4]** If *G = Vn* is an arithmetic graph, where , ai ≥ 1 for *i* ∈ {1, 2, . . . , *r*} , then the number of vertices of *G* can be calculated using the formula .

**Example 2.4.** Consider the graph *G* = *V*210 where 210 = 2 × 3 × 5×7 given in Figure 2.1. By Theorem 2.3 the number of vertices is |*V*210| = (1 + 1)(1 + 1)(1 + 1)(1 + 1) − 1 = 15.

**Theorem 2.5.** Let *Vn* be an arithmetic graph with , for any vertex where B ⊆ {1, 2, 3, . . . r}, 1 ≤ αi ≤ *a*i , and B ′ ⊆ B, one has deg(*u*) = (|B − B′ | +) − δ1|B| − ( δ1B′, where *B* is the number of distinct prime factors in a chosen vertex *u*, *B*′ is the number of prime factors having power 1 in chosen vertex *u*, empty summation equals zero and δ is the Kronecker’s delta function defined by δij =

**Note 2.6.** For an arithmetic graph the vertices are divisors of *n*, hence the degree of all types of vertices say primes, power of primes, product of primes, product of powers of primes can be calculated using the formulae given in Theorem 2.5.

**Theorem 2.7.** Let *G* = *V*n be an arithmetic graph where , then

∆(G) =

δ(G) = .

**Theorem 2.8.[4]** Let *G* = *V*n be an arithmetic graph where , then (i) ∆(G) = 2*r*−1

(ii)δ(G) =

**Theorem 2.9.** Let *G* be a *Vn* arithmetic graph, where, such that at least one of *i* ∈ {1, 2, . . ., *r*} does not equal one. Then,

(i) ∆(G) = *aj* where *a*j is the maximum exponent of *p*i, *i* ∈ {1, 2, . . ., r}

(ii)δ(G) = r.

**Results 2.10**.

1) It is identified that given arithmetic graph *G* = *V*n, ; ai ≥ 1 are bipartite.

2) The size of the arithmetic graph *G* = *V*n, where a1, a2 ≥ 1 is determined using the formula = 4 .

3)The diameter of an arithmetic graph is diam(G) ≤ 3 and its radius is radius(G) ≤ 2 .

4)Arithmetic graph *G = Vn*, is Hamiltonian, if , 3 ≤ r ≤ 6 .

5)All arithmetic graphs *G = Vn* is not a Eulerian graph.

**3.Connectivity Number of an Arithmetic graph *G = V*n**

Connectedness plays an essential role in graph theory, the graph representing a communication network needs to be connected for communications to be possible between all nodes(vertices). Numerous networks such as transport networks, road networks, electrical networks, distributed computing, block chain network, telecommunication systems or networks of servers can be modelled by a graph. Many researchers are made to determine how well a network is connected or can be splitted for sake of effectiveness. Two classical measures that indicate how the graph *G* is reliable are the edge-connectivity κ′(*G*) and the vertex-connectivity or simply the connectivity κ(*G*) of *G*. The connectivity of G, written κ(*G*), is the minimum order of a vertex set S ⊂ V (G) such that *G − S* is disconnected or has only one vertex.

**Definition 3.1.** The **connectivity or vertex connectivity κ(*G*)** is the number of vertices of a minimal vertex cut. A graph is called *k* − connected or *k* − vertex connected if its vertex connectivity is *k* or greater. Any graph *G* is said to be *k*-connected if it contains at least *k* vertices, but does not contain a set of *k* −1 vertices whose removal disconnects the graph and κ(*G*) is defined as the largest *k* such that *G* is *k* − connected. Thus κ(*G*) = 0 if *G* is either trivial or disconnected. All non-trivial connected graphs are 1 − connected.

**Definition 3.2.** An edge cut of *G* is a set of edges whose removal renders the graph *G* disconnected. The **edge-connectivity κ ′ (*G*)** is the size of a smallest edge cut. A graph is called *k* − edge − connected if its edge connectivity is *k* or greater. All non-trivial graphs are one edge connected.

**Theorem 3.3.** For an arithmetic graph *G* = *V*n, where and are distinct primes, then

**Theorem 3.4.** For an arithmetic graph *G = Vn*, where *p*i , i = 1, 2, . . . , *r, r* > 2 are distinct primes and *a*i = 1 for all *i* = 1, 2, . . . , r, κ(*Vn*) = κ ′ (*Vn*) = *r*

**Theorem 3.5.** For an arithmetic graph *G = Vn*, where are distinct primes and ai′ s ≥ 1 for all *i* = 1, 2, 3, . . . , *r* and *r* > 2 then κ() = κ ′ () = *r*.

**Proof**. We prove the theorem by considering the following four cases

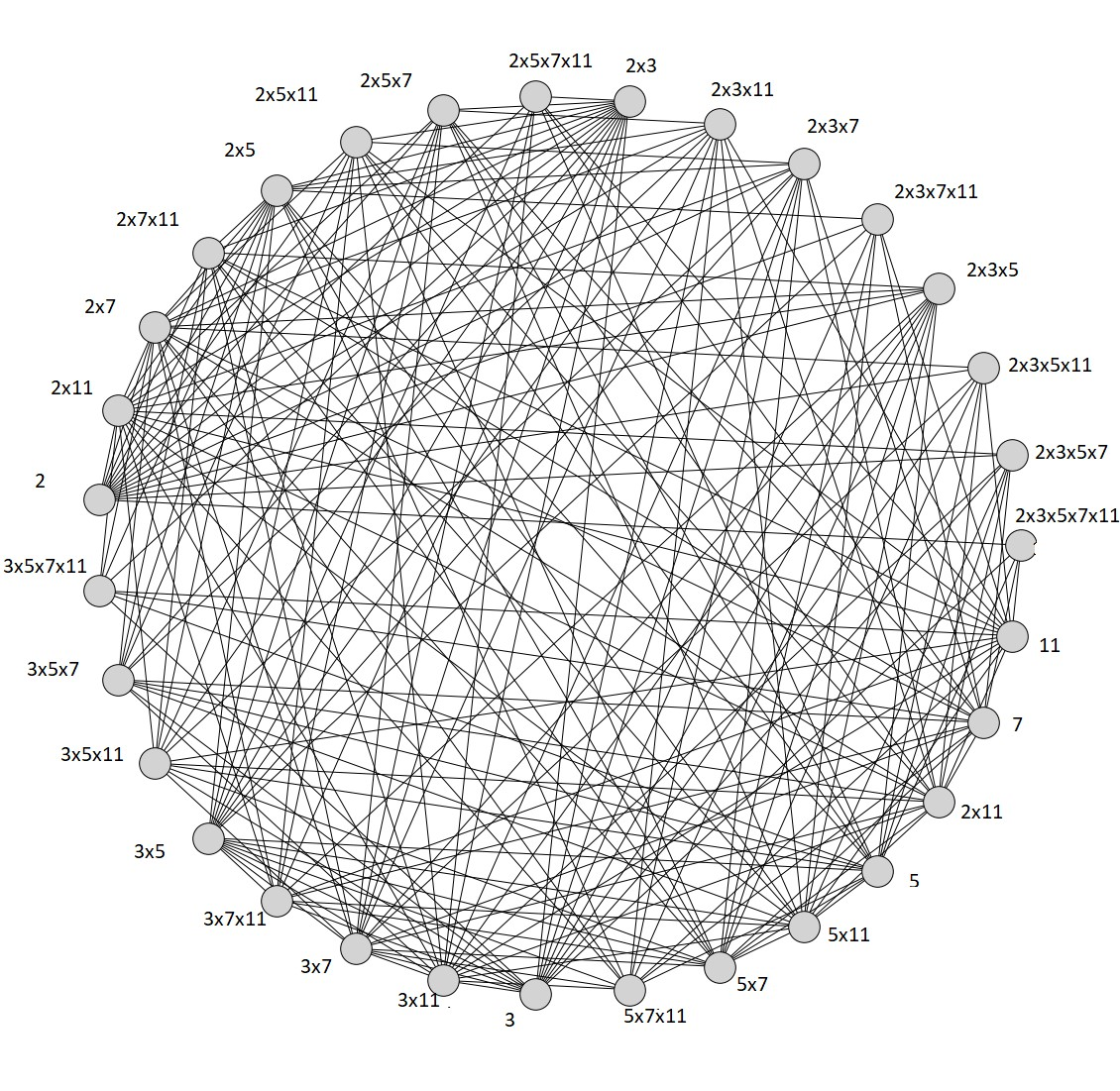
**Case (i)** All the *a*i’s, i = 1, 2, 3, . . . r is equal to one. By Theorem3.2., the result follows.

**Case (ii)** Some of the *a*i’s are equal to one and the others are greater than 1. Consider the vertex set of *V*n as *V (Vn*) = {*p1, p2, . . . , pr, p1 × p2*, . . . , }. Let the last vertex be say *v*1, where *ai* ’s are the maximum powers of the given distinct primes. By the definition of an arithmetic graph, we see that the only vertices which are adjacent to *v*1 are *p1, p2, . . ., pr*. Hence d(*v*1) = *r*. Also, the minimum degree of *V*n occurs at the vertex *v*1. That is, δ(*V*n) = *r* = *d(v1*). Hence, κ(*V*n) = κ ′ (*V*n) ≤ δ(*V*n) = *r*. But the removal of *r* vertices adjacent to *v*1 makes the graph disconnected. Hence, we obtain the result κ(*V*n) = *r*. The edge connectivity κ ′ (*V*n) = *r* is same as Theorem 3.2.

**Case(iii)** All the *ai* ’s are equal and greater than 1. Here also consider the last vertex of *V* (*V*n), say where the *a*i’s are the maximum power of given distinct primes. By the definition of an arithmetic graph, it is clear that *p1, p2, . . . , pr* are the only vertices which are adjacent to the vertex . The remaining proof is similar to case (ii).

**Case (iv)** All the *a*i’s are distinct and greater than one. Consider the last vertex in the vertex set of *V*n, say where the *a*i’s are the maximum power of the given distinct primes. By the definition of an arithmetic graph, this vertex is adjacent to exactly r vertices namely *p1, p2, . . . , pr*. Suppose it is adjacent to any other vertex except *p*i then, it contradicts the definition of an arithmetic graph. The remaining proof is similar to case (ii).

**Example 3.6.** Consider an arithmetic graph *G = V*2310 where 2310 = 2 × 3 × 5 × 7 × 11, given in Figure 3.1. The set *S* = {2, 3, 5, 7, 11} is a minimum vertex cut so the cardinality of the set *S* is a connectivity number κ(*G*) which is equal to 5. Also, the removal of the set of edges say *F* = {2 × 3 × 5 × 7 × 11 2, 2 × 3 × 5 × 7 × 11 3, 2 × 3 × 5 × 7 × 11 5, 2 × 3 × 5 × 7 × 11 7, 2 × 3 × 5 × 7 × 11 11} makes the graph disconnected, since *F* is a minimum edge cut, cardinality of *F* is the edge connectivity number. Thus, we have κ′(*G*) = 5



**Figure 3.1: Arithmetic graph *G = V*2310**

**Remark 3.7.** The arithmetic graph *Vn* is a maximally connected graph

**4. The connectivity number of complement of an arithmetic graph**

In this section, we identified the connectivity number for complement of an arithmetic graph *G = Vn,* where *n* is a product of two primes, product of powers of two primes, product of *r* primes, product of powers of *r* primes.

**Definition 4.1** Let *G* be a graph, The complement of a graph *G* is the graph with vertex set *V* () such that two vertices are adjacent in if and only if they are not adjacent in *G*

**Theorem 4.2** For an arithmetic graph *G = Vn*, where *a*j > 1;, κ () = *aj* – 1

**Theorem 4.3** For an arithmetic graph *G = Vn*, where *a*j ≥ *a*i ≥ 2, κ () = *ai + aj* – 1

**Example 4.5** Consider the graph *G = V36*, 36 = 22 × 32 given in Figure 4.1. The set S1 = {2, 3} is a minimum vertex cut, and the Figure 4.2 shows the complement of an arithmetic graph *V*36 and the set *S*2 = {22 , 3, 32} is the minimum vertex cut of *G*. Thus κ () = 3 which satisfies the value of *ai + aj* − 1 = 2 + 2 − 1 = 3

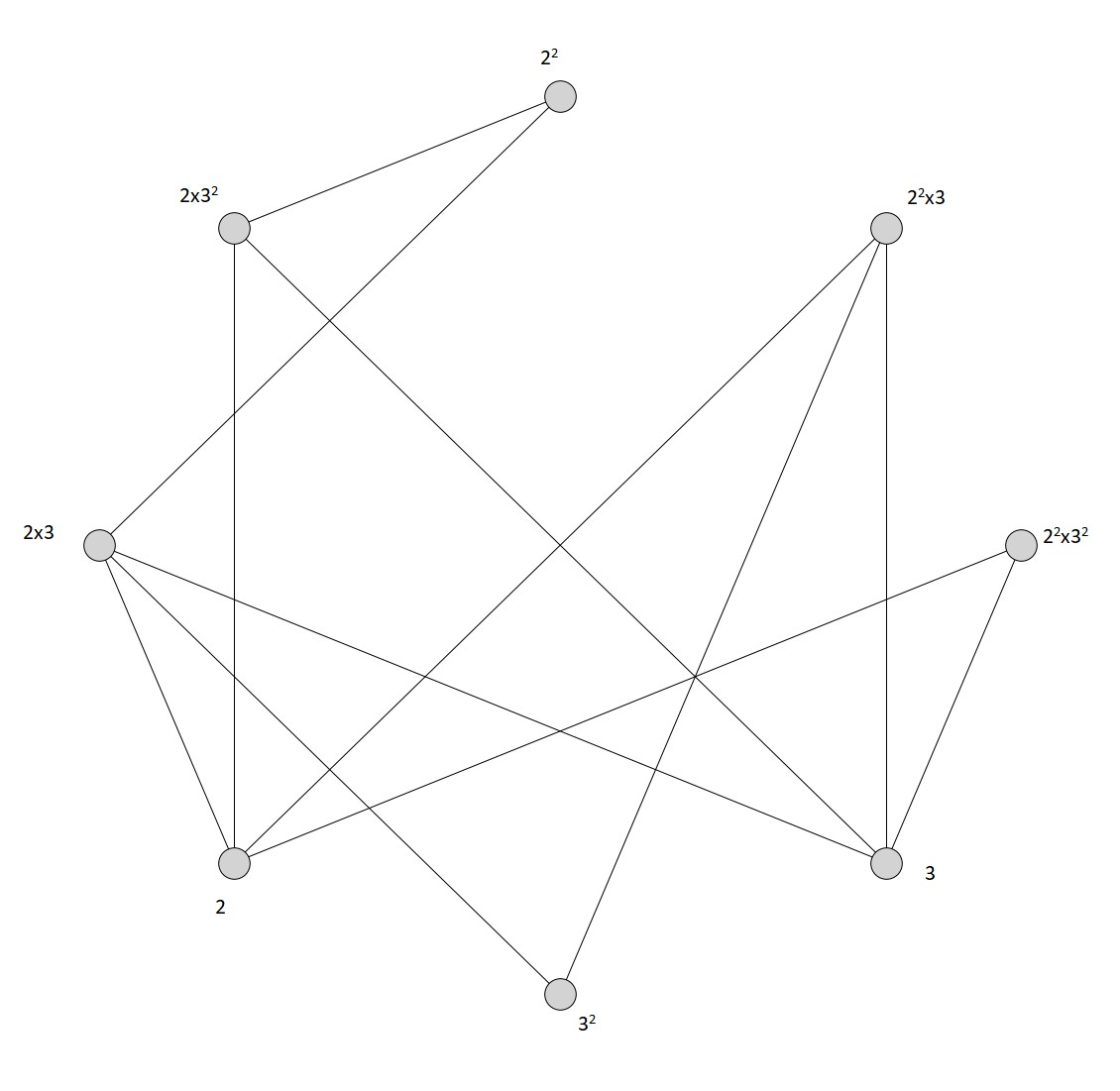


Figure 4.1: Arithmetic Graph *G=V*36

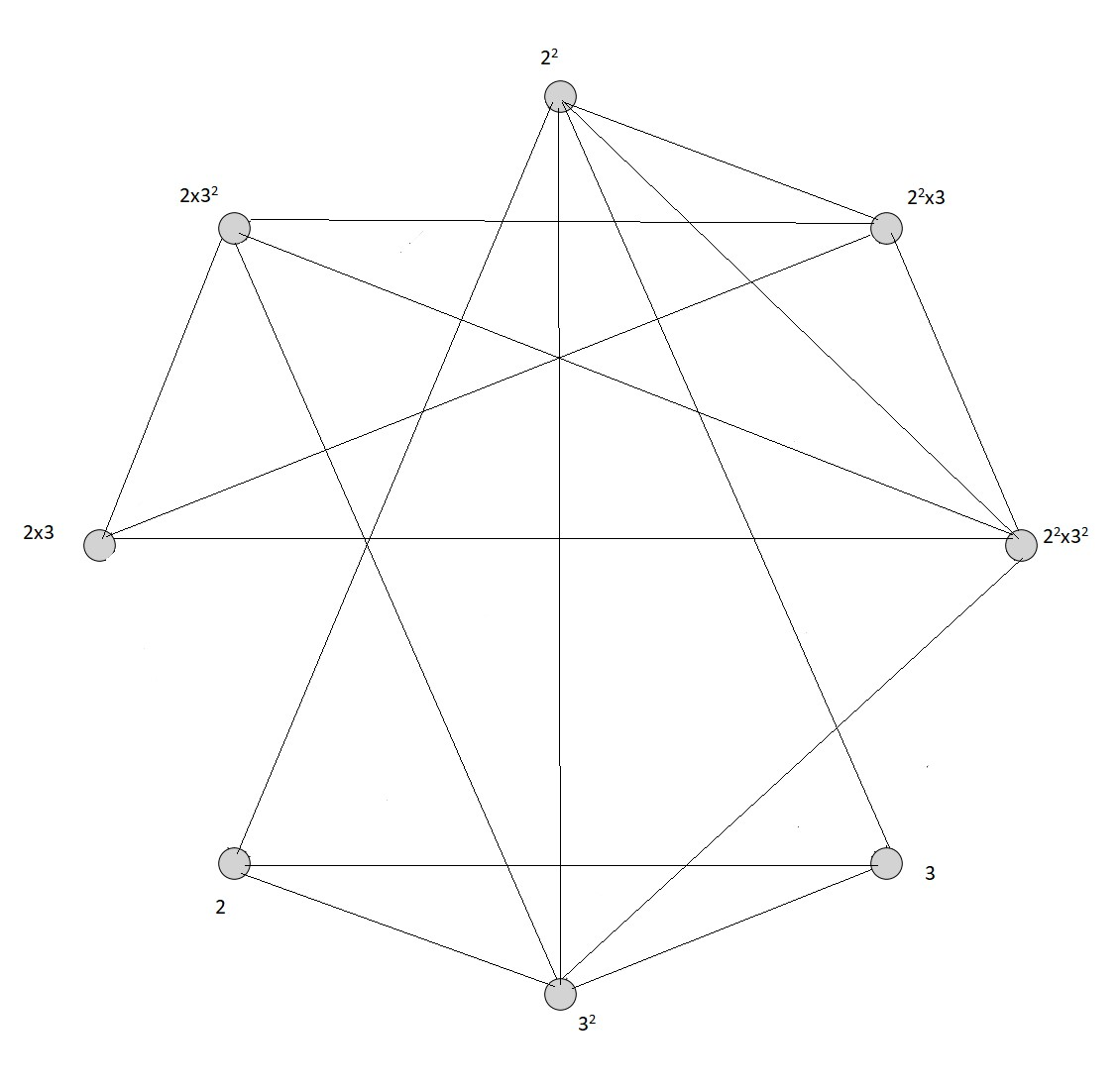


Figure 4.2: Arithmetic graph  *=*

**Theorem 4.6** For an arithmetic graph *G = Vn*, , *r* > 2, κ () = (2*r* − 4) /2.

**Proof**. Let *G = Vn* be an arithmetic graph. The vertex set *V (G) = V* () = {*, , . . . , , p1×p2, . . . , p1×pr, . . . , pr−1×pr, p1×p2×p3, . . . ,* }. By Theorem 2.2.4, the maximum degree of the graph *G* is ∆(*G*) = 2*r*−1 . The vertices which are of having maximum degree is *pi × pj*;i, j ∈ {1, 2,. . . , *r*}, . Hence the number of vertices having maximum degree is *rC*2. These *rC*2 vertices will have minimum degree in . Thus, there are *rC*2 set of minimum vertex cuts. Also, the minimum degree δ() = |*V (G*) − 1| − ∆ (*G*) = (2*r* − 4) /2. Thus *d (pi × pj )* = (2*r* − 4) /2. So we have κ( = (2*r* − 4) /2.

**Theorem 4.7** For an arithmetic graph *G* = *V*n, , *r* > 2 and at least one of *ai* , *i* ∈ {1, 2, . . . , *r*} does not equal to one, where *aj* is the maximum exponent of *pi*, i ∈ {1, 2, . . . , *r*}.

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